

Introduction

I set out to bring the world of Intel Hardware Transactional Memory (HTM) to Go.

To make this happen, I would need to present the following items:

- A library that made it easy to use HTM in Go
- Evidence that my primitives worked as intended and were special
- Some real-world gains in common data structures

Background

Intel has implemented transaction memory features under the name Intel TSX.

Intel TSX contains two different instruction extensions, HLE and RTM, which approach HTM from two different angles.

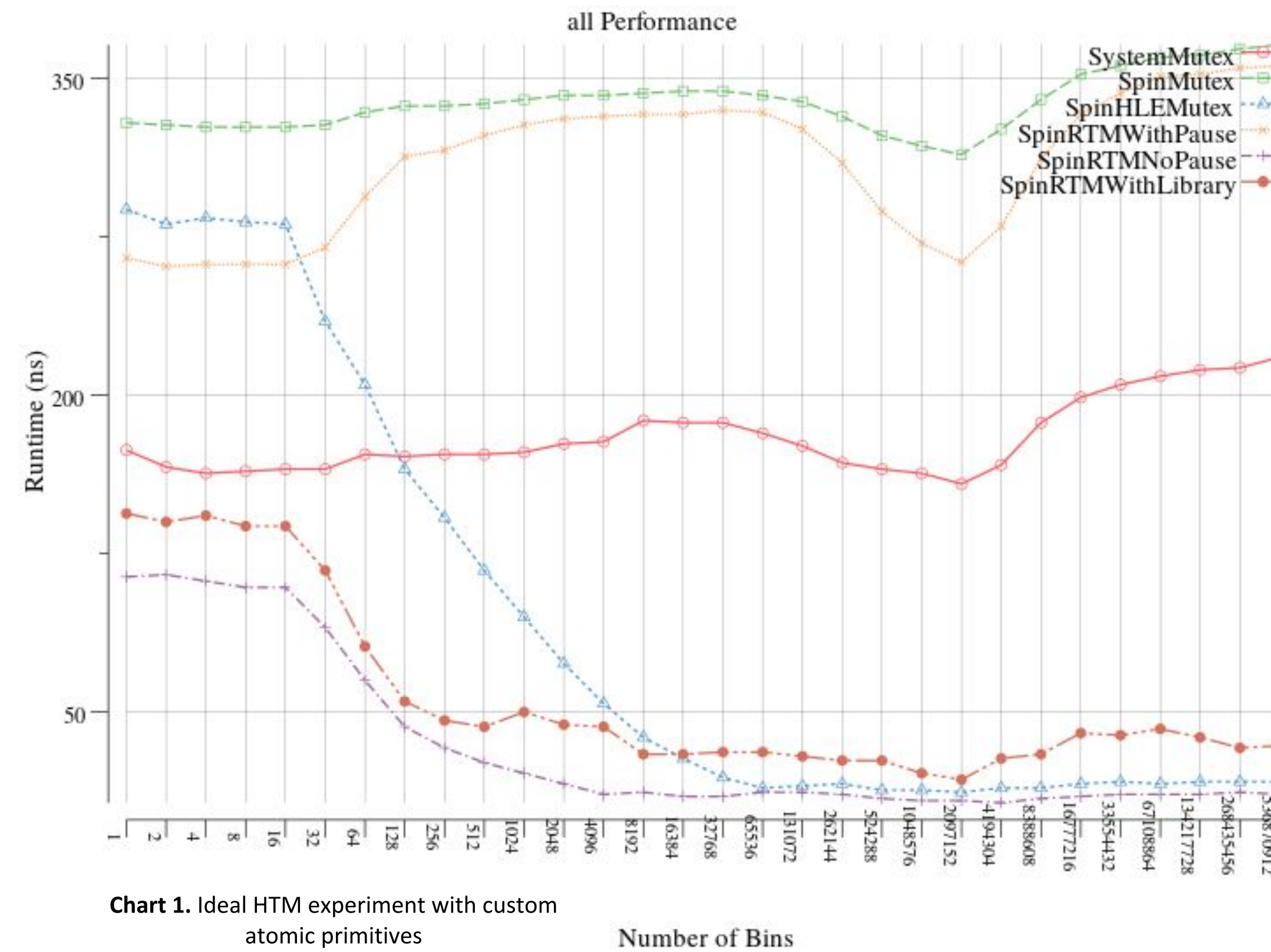
Simply put, **HLE** is designed as a drop-in replacement for current locks, whereas **RTM** allows low level control of transactions.

Results

I have implemented highly tuned Go primitives to make use of both HLE and RTM.

In the ideal HTM experiment, the new primitives are validated and exceed the expected performance. The hash-map experiment (and AVL tree experiment) stand to demonstrate the gains from using these primitives in real world applications.

A simple take away is that HTM can massively speed up reads or modifies, but often fails during insertions, due to global data structure tampering.



```
var lock SpinHLEMutex
lock.Lock()
// Action to be done transactionally
count := m["word1"]
m["word1"] = count + 1
lock.Unlock()
```

Figure 1. Example of using the HLE lock replacement

```
c := NewRTMContextDefault()
c.Atomic(func() {
    // Action to be done transactionally
    count := m["word1"]
    m["word1"] = count + 1
})
```

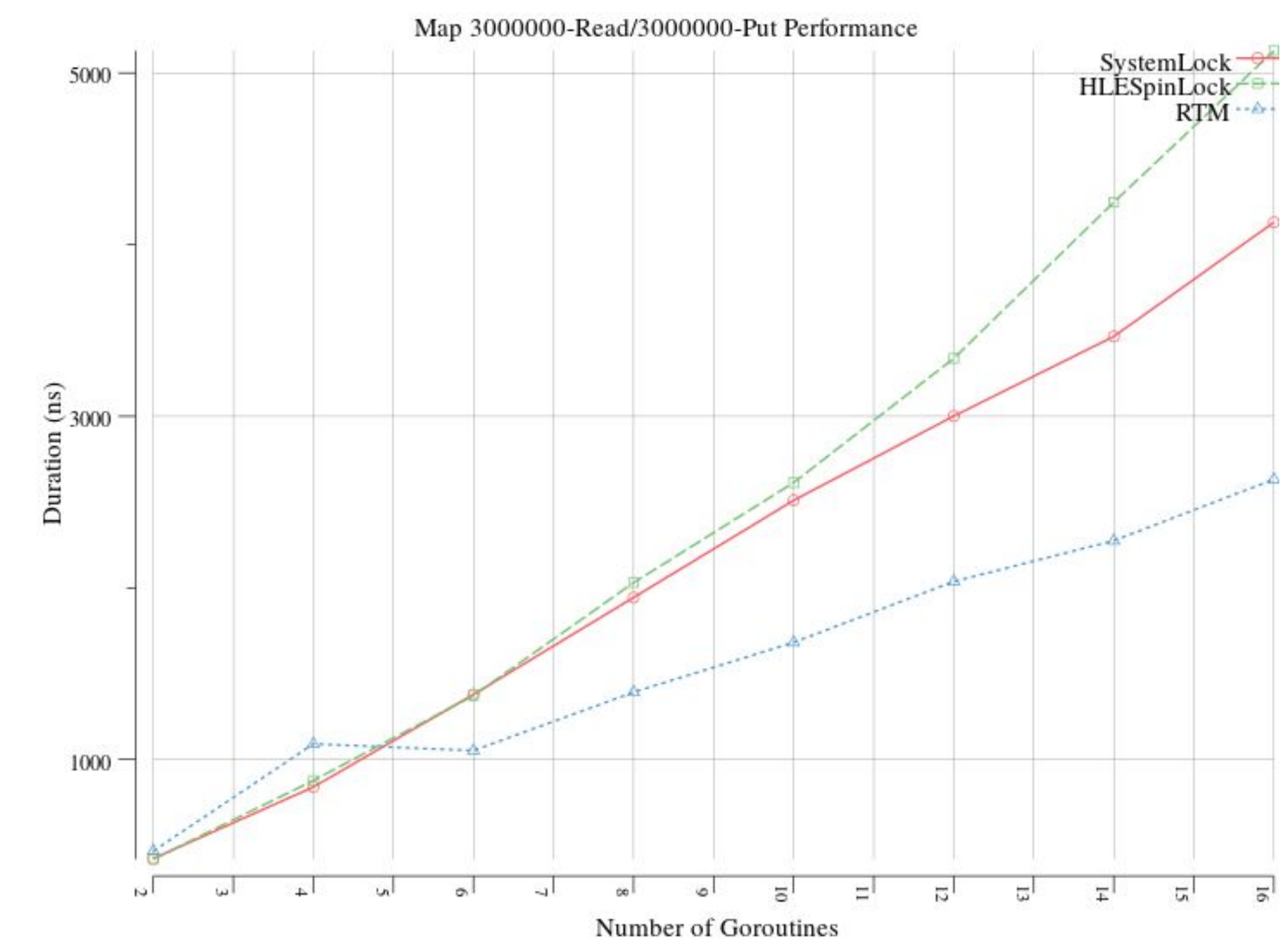
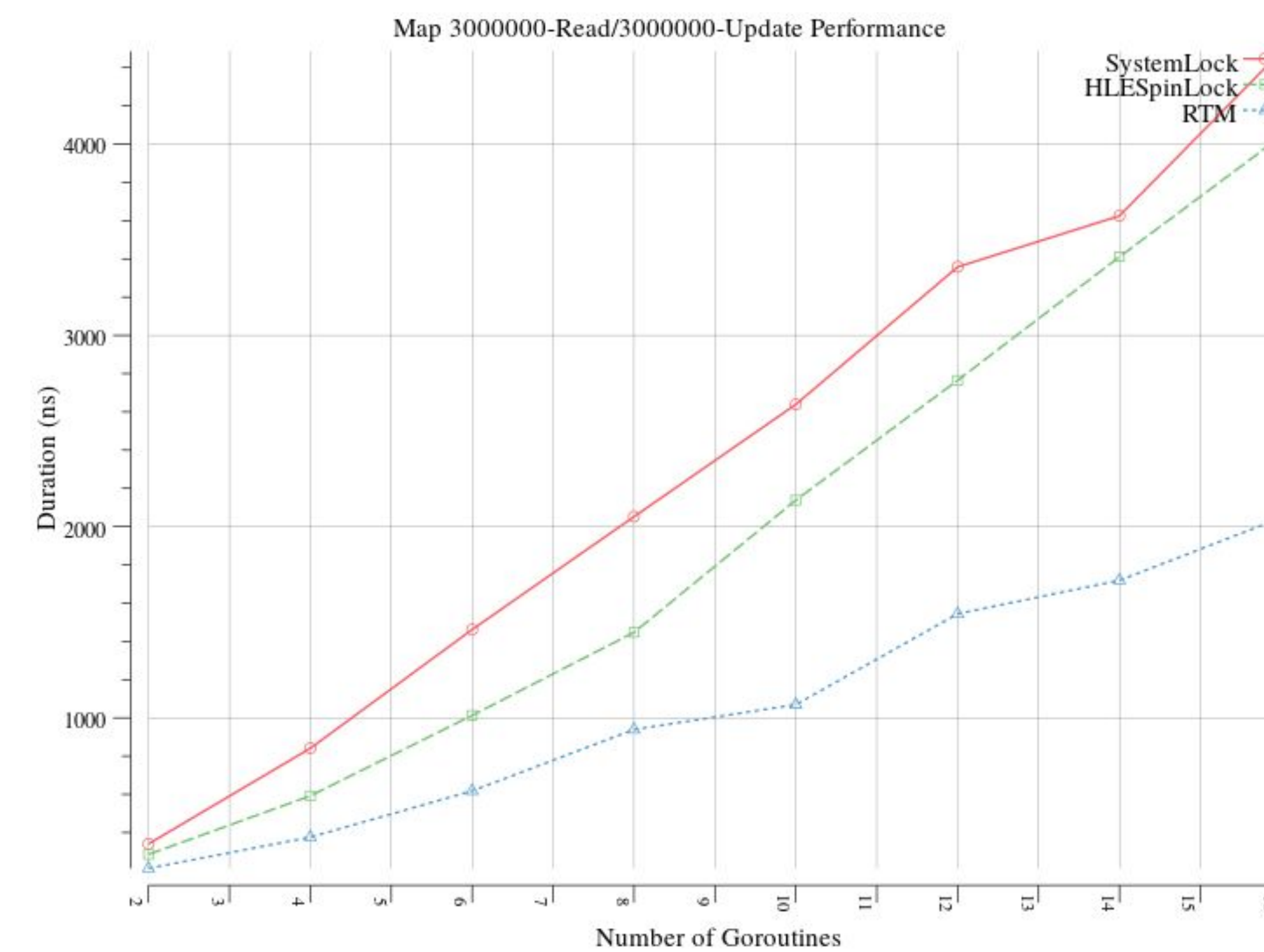
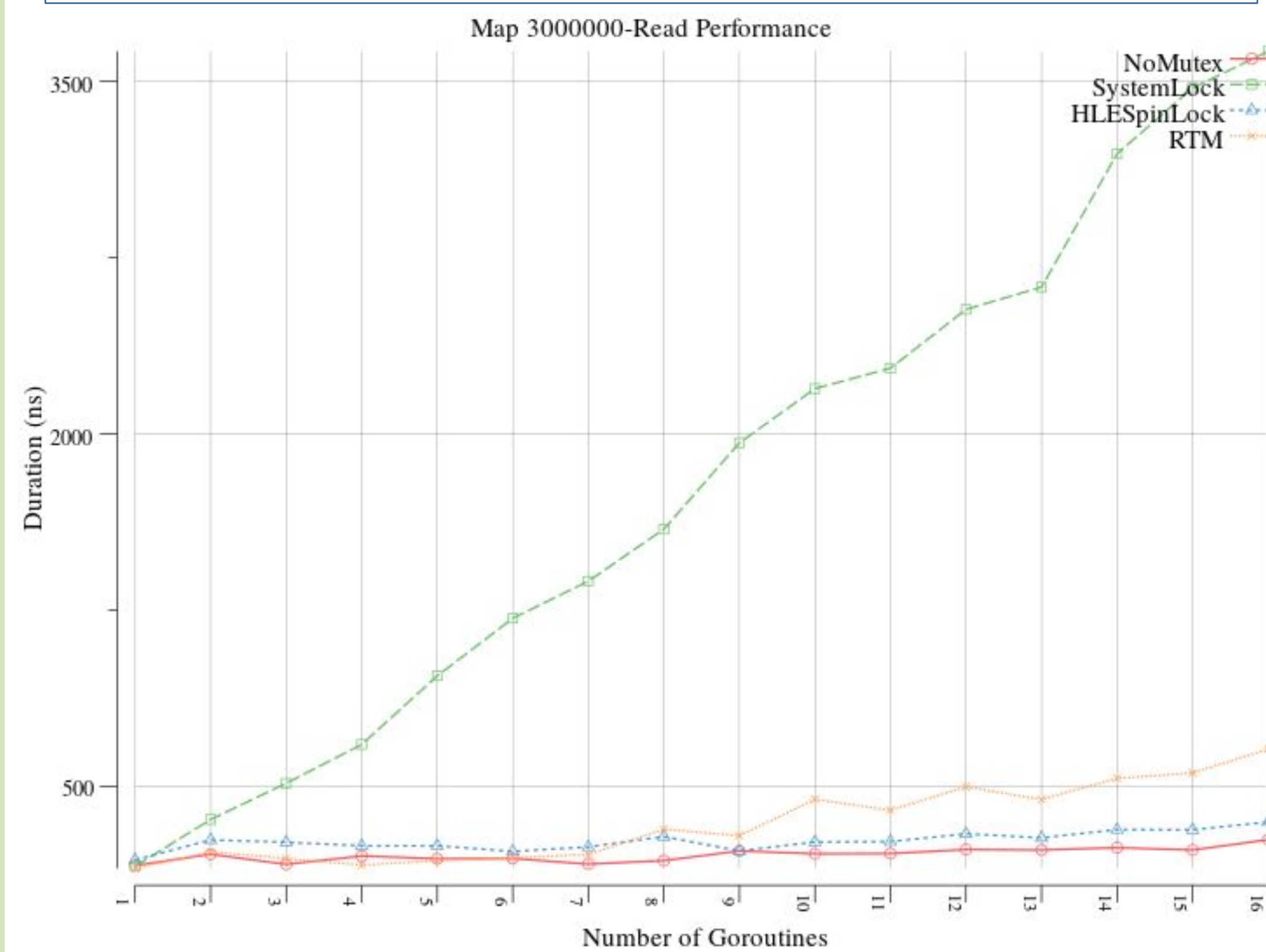
Figure 2. Example of using the RTM transaction context primitive.

```
c := NewLockedContext(new(SpinHLEMutex))
c.Atomic(func() {
    // Action to be done transactionally
    count := m["word1"]
    m["word1"] = count + 1
})
```

Figure 3. Example of using the HLE transaction context primitive. This allows interoperability between RTM, HLE, and normal system locks.

```
114 // Note: Argument attempts must be greater than 0
115 // func HLESpinCountLock(val, attempts *int32)
116 TEXT ·HLESpinCountLock(SB),NOPTR|NOSPLIT,$0
117     MOVQ val+0(FP), CX
118     // Load attempt counter in DX
119     MOVQ attempts+8(FP), R8
120     MOVL (R8), DX
121 tryread:
122     MOVL (CX), BX
123     TESTL BX, BX
124     JE tryacquire
125     PAUSE
126     DECL DX
127     // If DX != 0, abort
128     JNE tryread
129     JMP abort
130 tryacquire:
131     MOVL $1, AX
132     XACQUIRE
133     XCHGL AX, (CX)
134     TESTL AX, AX
135     JNE tryread
136 abort:
137     // Write back attempt counter
138     MOVL DX, (R8)
139     RET
```

Figure 4. Implementation of the HLE spin lock in Go ASM



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Project Page

<https://github.com/linux4life798/safetyfast>